

Lecture Chapter 6 Recursion as a Problem Solving Technique

Backtracking

1. Select, i.e., guess, a path of steps that could possibly lead to a solution
2. If the path leads to a dead end then retrace steps in the reverse order
3. Select a new sequence of steps that could possibly lead to a solution
4. If need be, go to #2 above

Eight Queens Problem

- **Chessboard**
 - 64 squares -- 8 rows x 8 columns
 - Queen can attack any other piece
 - within its row
 - within its column
 - along any diagonal
 - place eight queens on the board such that
no queen can attack any other queen,
 - number of ways to arrange 8 queens on a 65 square board is $c(64, 8)$
 - $c(64, 8) > 4$ trillion
thus each row & column contains exactly one queen
 - attacks along rows or columns are not possible
 - number of attacks to be checked along diagonals is $8! = 40,320$

○ Strategy 1

- Placing queens on the board in the following sequence yields a strategy which fails after five queens since column six is totally blocked

Q ₁	•	•	•	•	•		
•	•	•	Q ₄	•	•		
•	Q ₂	•	•	•	•		
•	•	•	•	Q ₅	•		
•	•	Q ₃	•	•	•		
•	•	•	•	•	•		
•	•	•	•	•	•		
•	•	•	•	•	•		

- Backtrack by removing Q₅ and finding a new location on column five.
- Placing Q₅ in any cell on column five still eliminates any queens being placed in column six.
- Backtrack by removing Q₄ & Q₅ and finding new locations on both columns.

Q ₁	•	•	•	•	•		
•	•	•	•	Q ₅	•		
•	Q ₂	•	•	•	•		
•	•	•	•	•	•		
•	•	Q ₃	•	•	•		
•	•	•	•	•	•		
•	•	•	Q ₄	•	•		
•	•	•	•	•	•		

A solution to the Eight Queens Problem

Q ₁	•	•	•	•	•	•	•
•	•	•	•	•	•	Q ₇	•
•	•	•	•	Q ₅	•	•	•
•	•	•	•	•	•	•	Q ₈
•	Q ₂	•	•	•	•	•	•
•	•	•	Q ₄	•	•	•	•
•	•	•	•	•	Q ₆	•	•
•	•	Q ₃	•	•	•	•	•

Languages

- Set of all Java Programs = $\{ \text{all strings } w : w \text{ is a syntactically correct Java program} \}$

Java Compiler : program which determines

if a string w is a syntactically correct Java program

- Set of all English Sentences = $\{ \text{all strings } w : w \text{ is a syntactically correct English Sentences} \}$

Set of all English Grammar Rules : determines

if a sentence w is a syntactically correct English Sentence

- Set of all Algebraic Expressions = $\{ \text{all strings } w : w \text{ is a syntactically correct Algebraic Expression} \}$

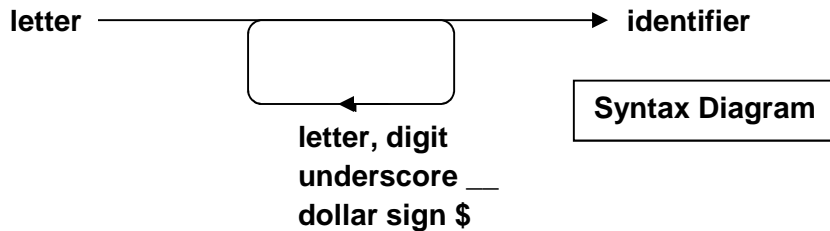
Set of all Algebraic Rules : determines

if a sentence w is a syntactically correct Algebraic Expression

- Grammar states the rules of a language
 - recursive rules – computer languages
 - non-recursive rules -- natural languages, e.g., English, Hungarian, etc.
- recognition algorithm – recursive algorithm based on the grammar
which determines whether a given string is in the grammar

Identifiers

- $x \mid y \leftrightarrow x \text{ or } y$
- $x y \leftrightarrow x \bullet y \leftrightarrow x \text{ concatenated with } y$
- $\langle \text{symbol} \rangle \rightarrow \text{any sequence of symbols defined by the grammar}$



- recursive grammar
 - $\langle \text{identifier} \rangle = \langle \text{letter} \rangle \mid \langle \text{identifier} \rangle \langle \text{letter} \rangle \mid \langle \text{identifier} \rangle \langle \text{digit} \rangle \mid \$ \langle \text{identifier} \rangle \mid _ \langle \text{identifier} \rangle$
 - $\langle \text{letter} \rangle = a \mid b \mid \dots \mid z \mid A \mid B \mid \dots \mid Z$
 - $\langle \text{digit} \rangle = 0 \mid 1 \mid \dots \mid 9$
- recognition grammar
 - $\text{length}(w) == 1 \rightarrow w \text{ is an identifier if } w \text{ is a letter}$
 - $\text{length}(w) \geq 1 \rightarrow w \text{ is an identifier if}$
 1. last character of w is either a letter or a digit
 2. w minus the last character is an identifier

```
isId(in w: string) : boolean
if ( length(w) == 1 )
{
    if ( w is a letter ) return true
    else return false
}
else if ( last character of w is either a letter or a digit )
{
    return isId( w minus the last character )
}
else return false
```

Strings $A^n B^n$

- n consecutive A's followed by n consecutive B's
- AAAAABBBBB

$L = \{ w : w \text{ is of the form } A^n B^n \text{ for some } n \geq 0 \}$

- Grammar <legal-word > = empty string | A < legal-word > B
- Recognition algorithm

isAnBn(in w : string) : Boolean

```
{
  if ( length(w) == 0 ) return true
  else if ( w begins with A and ends with B )
  {
    return isAnBn( w minus first & last characters )
  }
  else return false
}
```

Algebraic Expressions

Binary Operators: +, -, *, /

Operands: single letter only

- Infix, Prefix, Postfix Expressions
 - Infix
 - operand₁ operator operand₂ e.g., x + y
 - associative rules
 - precedence rules
 - use of parentheses
 - Prefix
 - operator operand₁ operand₂ e.g., + x y
 - infix a + (b * c) → prefix + a * b c
 - infix (a + b) * c → prefix * + a b c
 - Postfix
 - operand₁ operand₂ operator e.g., x y +
 - infix a + (b * c) → postfix a b c * +
 - infix (a + b) * c → postfix a b + c *

➤ Conversion Infix to Prefix &/or Postfix

- Infix → fully parenthesized infix

$$a + b * c \rightarrow ((a + b) * c)$$

- Fully Parenthesized Infix → Prefix

- Move each operator to the position marked by its **open** parenthesis

$$((a + b) * c) \rightarrow ((a \quad b) c) \rightarrow * + a b c$$

* +

- Fully Parenthesized Infix → Postfix

- Move each operator to the position marked by its **closed** parenthesis

$$((a + b) * c) \rightarrow ((a \quad b) c) \rightarrow a b + c *$$

+ *

Prefix & Postfix Expressions do not require

- associative rules
- precedence rules
- use of parentheses

to avoid ambiguity

- Prefix Grammar

<prefix> = <identifier> | <operator> <prefix><prefix>

<operator> = + | - | * | /

<identifier> = a | b | ... | z

To be Expanded Later in the Semester

- Postfix Grammar

<postfix> = <identifier> | <postfix> <postfix> <operator>

<operator> = + | - | * | /

<identifier> = a | b | ... | z

To be Expanded Later in the Semester